

Heap Taichi: Exploiting Memory Allocation Granularity in Heap-Spraying Attacks

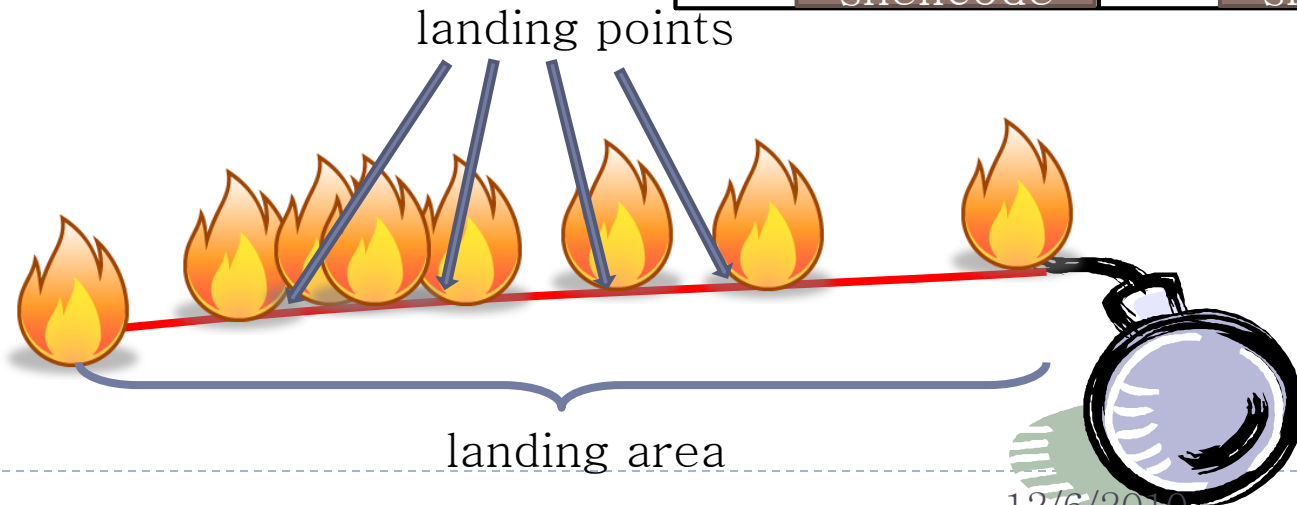
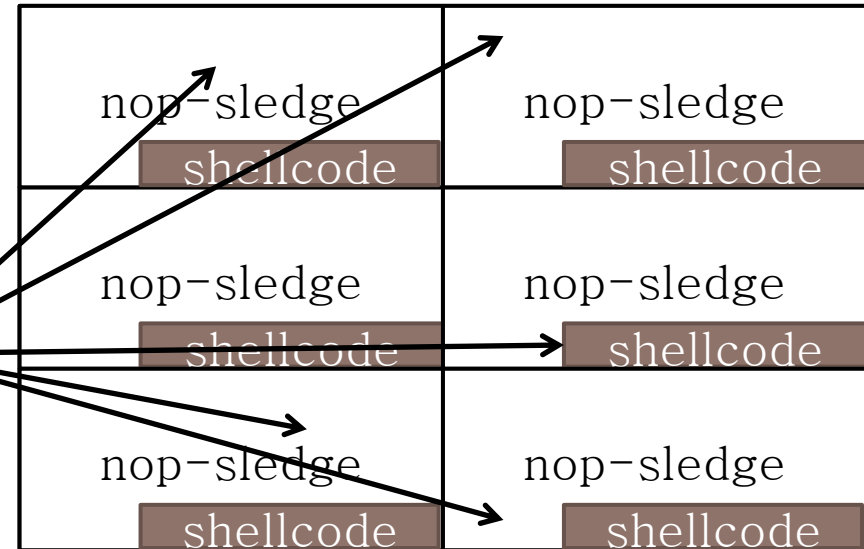
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A Typical Heap-Spraying Attack

```
var cc = unescape("%u0C0C");  
var sc = unescape("%u785C%u3334...");  
while (cc.length * 2 < 0x100000)  
    cc+=cc;  
cc=cc.substring(sc.length + 0x38,  
                0x100000/2);  
var m = new Array();  
for(i = 0; i < 200; i++)  
    m[i] = cc + sc;
```

PC

Heap Area



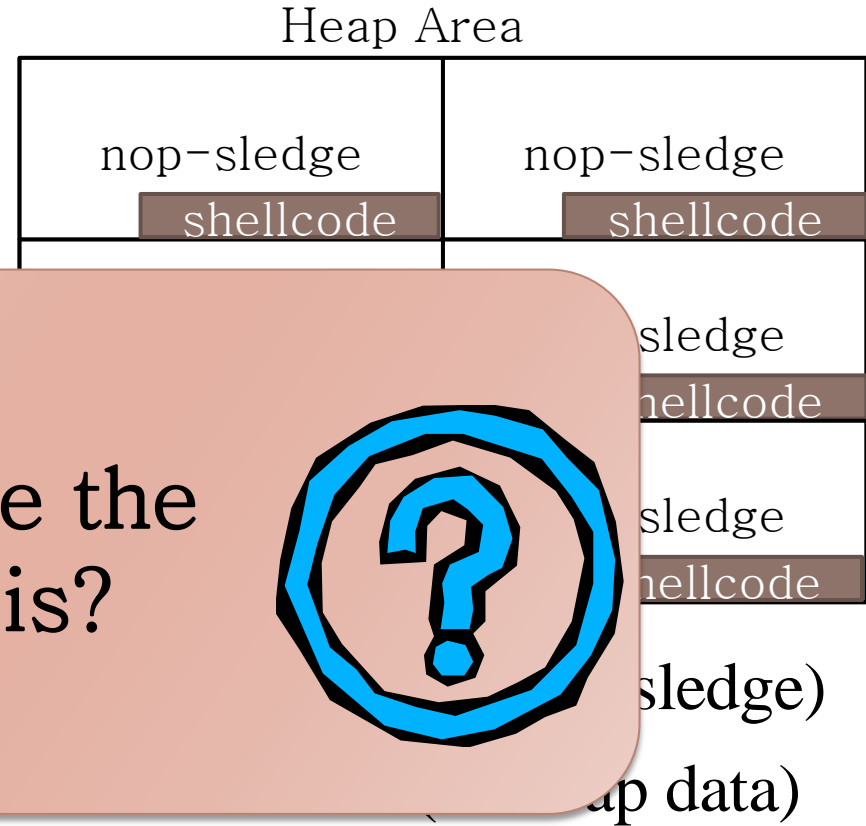
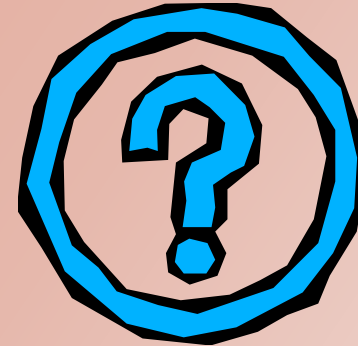
Existing Solutions

Measure how much of the heap can be used as nop-sledge

Because of randomization (ASLR), the large amount of nop instructions is required.

A direct d
measure h
can be use
flows to a

But how reliable the randomization is?



Surface A

Normalized

If the NSA > threshold, generate an alarm.

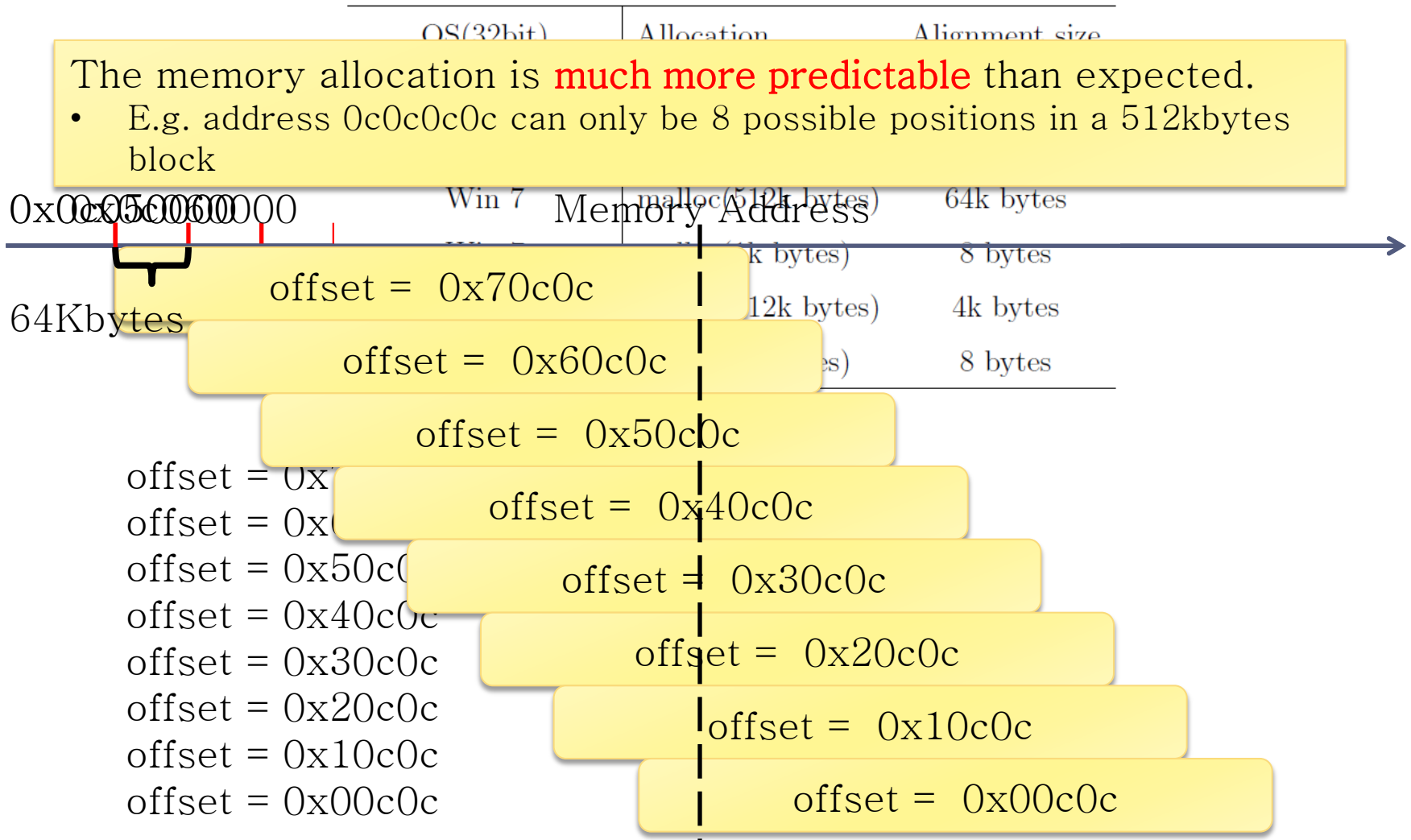
Nozzle [USENIX'09]

Observation

Memory Allocation Granularity

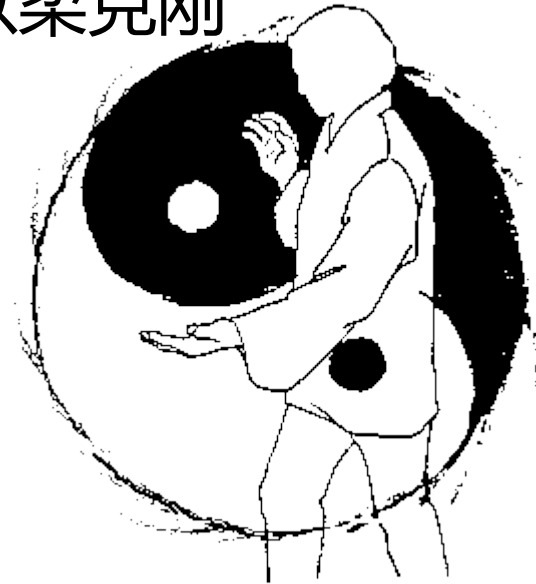
The memory allocation is **much more predictable** than expected.

- E.g. address 0c0c0c0c can only be 8 possible positions in a 512kbytes block



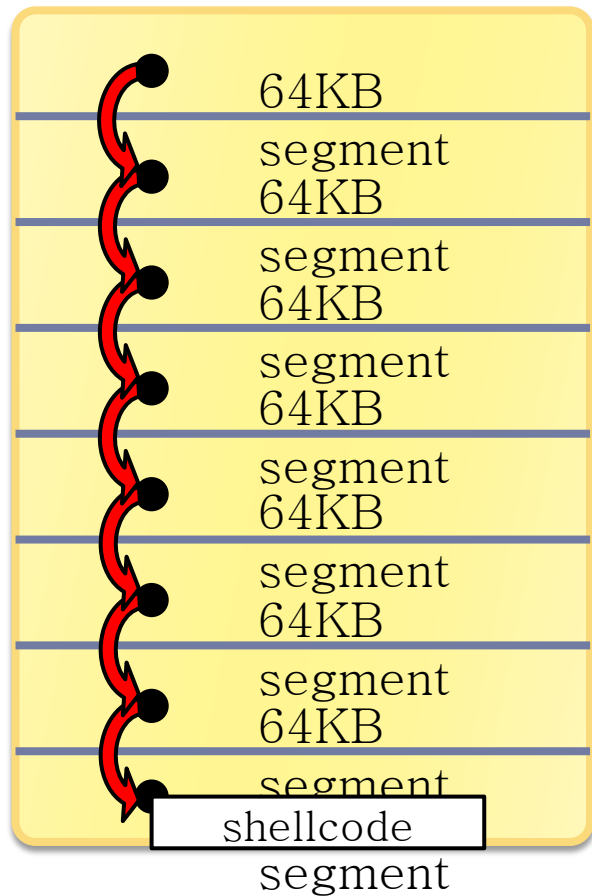
Proof of Concept: Heap Taichi

- ▶ Memory allocation granularity makes defense mechanisms less flexible
- ▶ We discuss a few transformations to reduce heap-spraying attacks' footprint without losing its ...
- ▶ Countering force with agility 以柔克刚

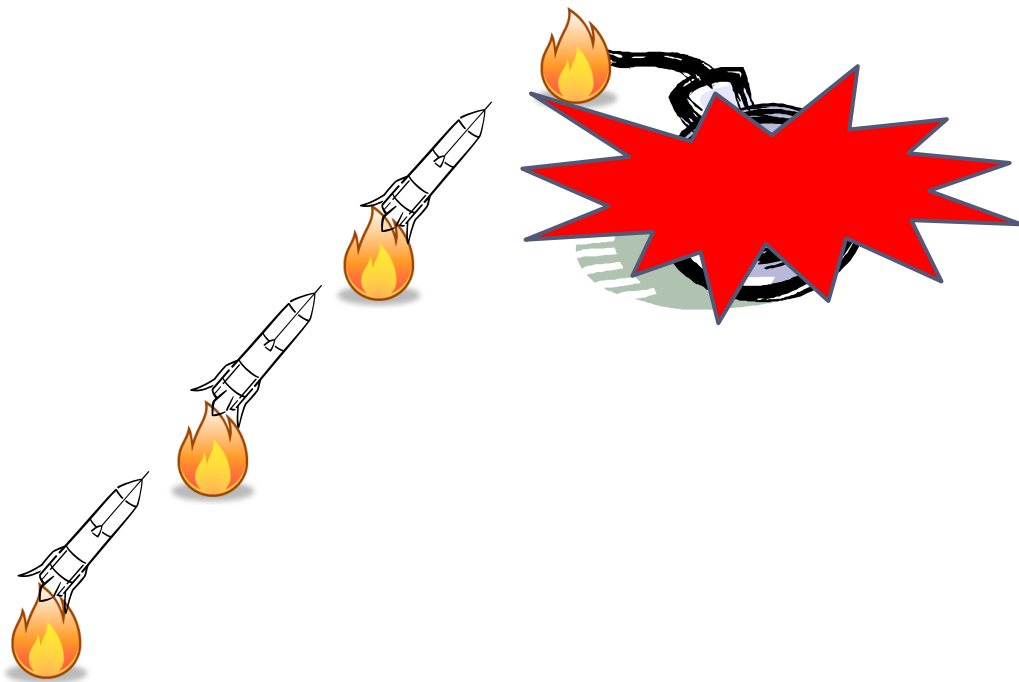


Basic Structure of TypeA (Passing Flower)

alignment size = 64Kbytes
blocksize = 512Kbytes

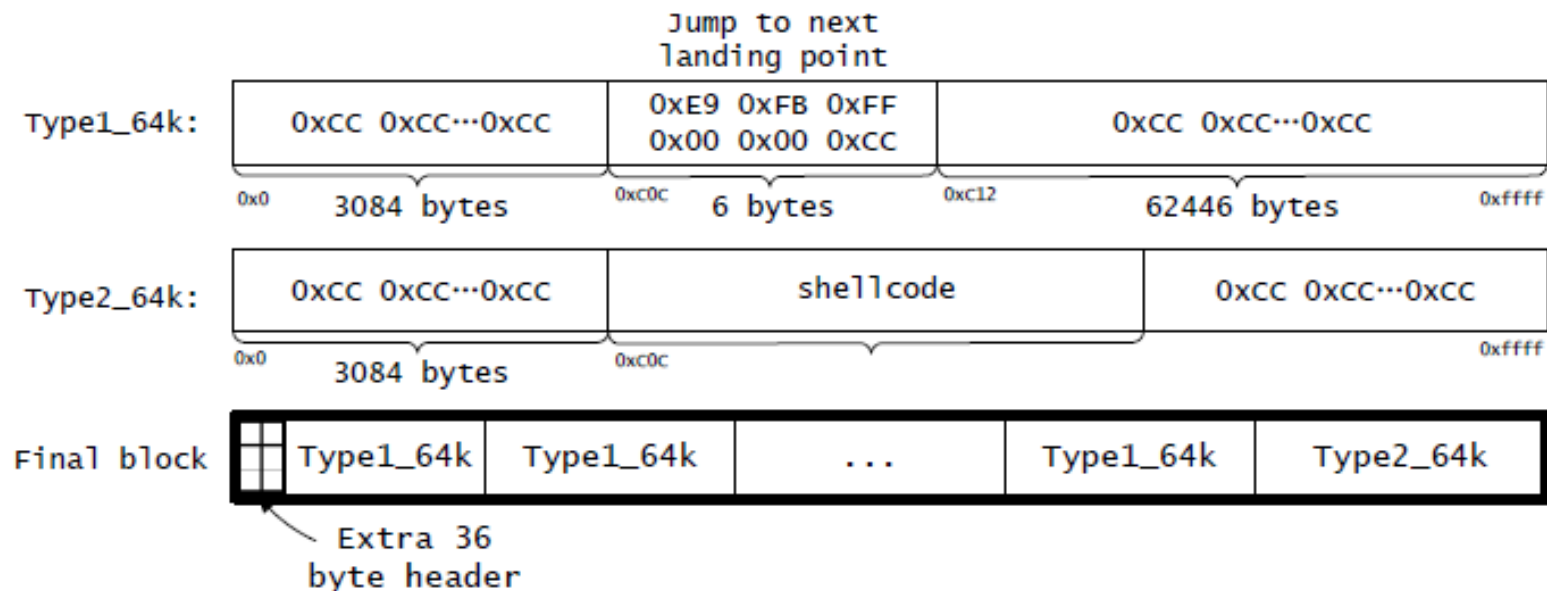


NSA(sledge) $\approx 7 / 512K \approx 0\%$



The Structure of TypeA Attack (Detail)

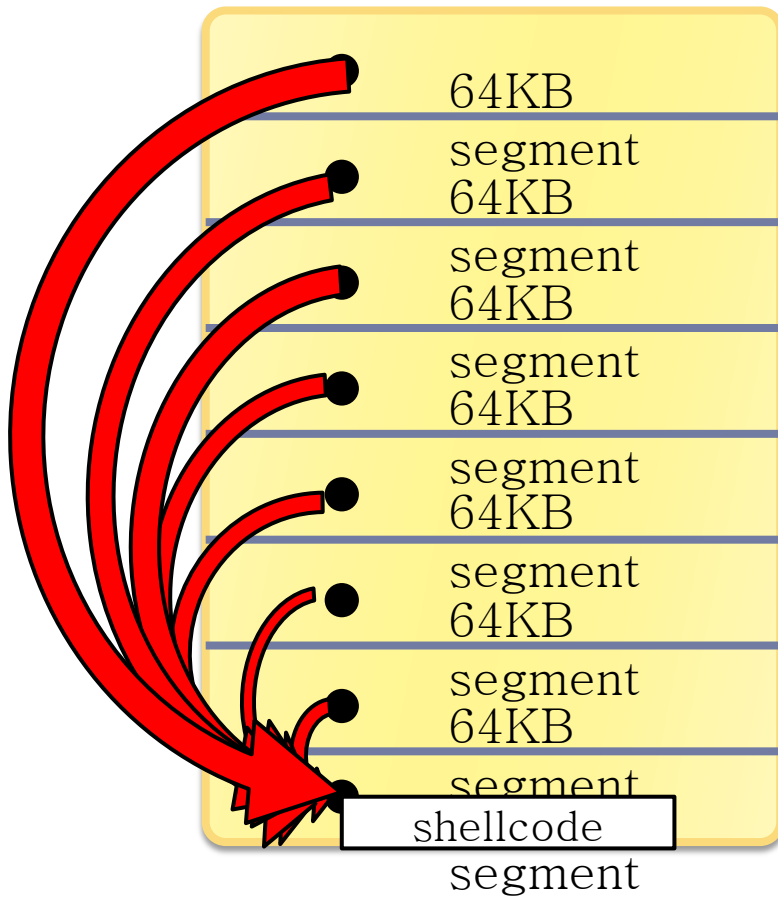
- ▶ A sample JavaScript code creating Type A heap objects



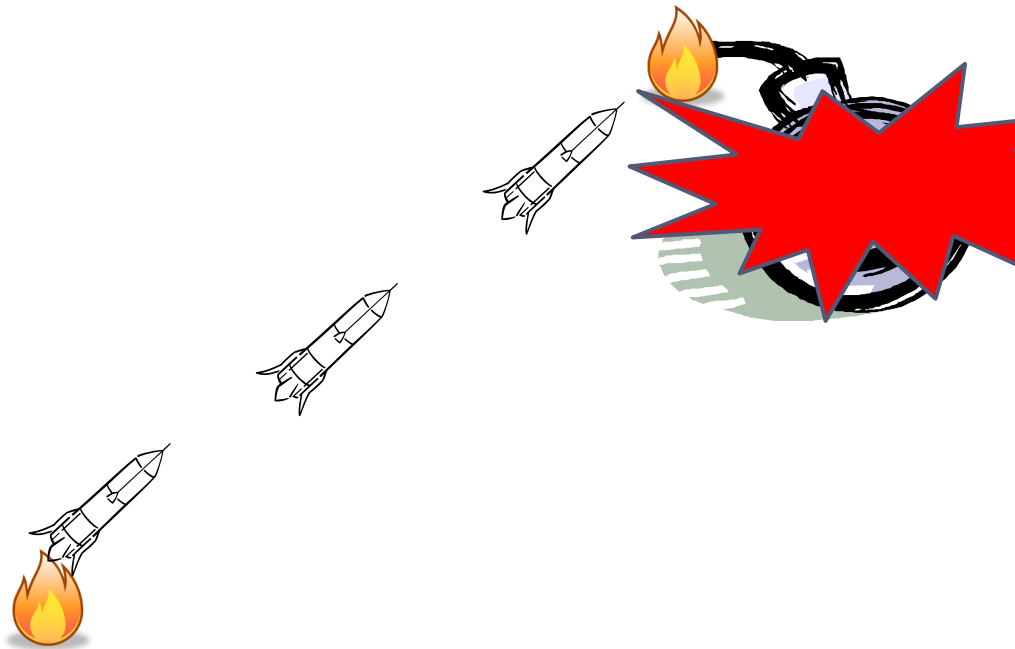
- ▶ The size of each heap memory block is 1Mbytes.
- ▶ Try to allocate 200 heap memory blocks to cover address 0x0c0c0c0c

Basic Structure of TypeB (Jumping together)

alignment size = 64Kbytes
blocksize = 512Kbytes



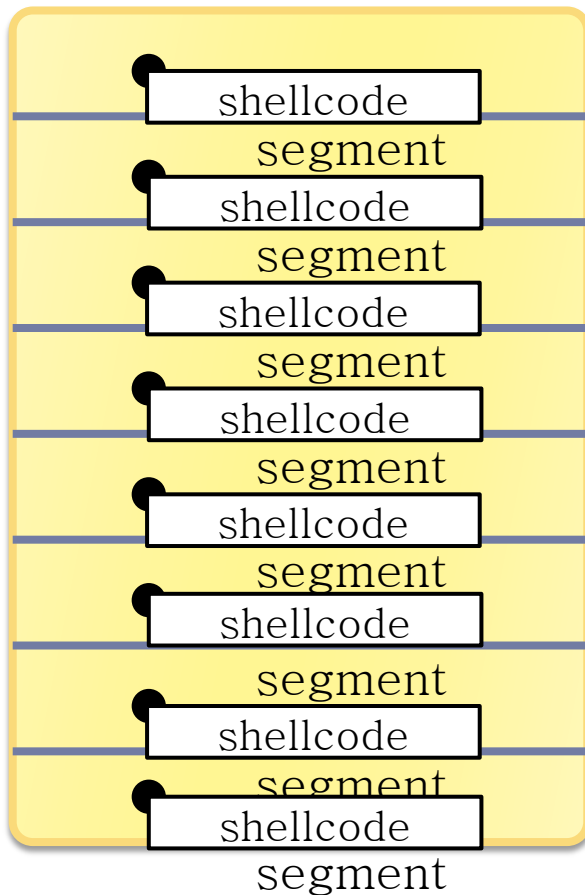
NSA(sledge) $\approx 7 / 512K \approx 0\%$



Basic Structure of TypeC (Returning Home)

alignment size = 64Kbytes

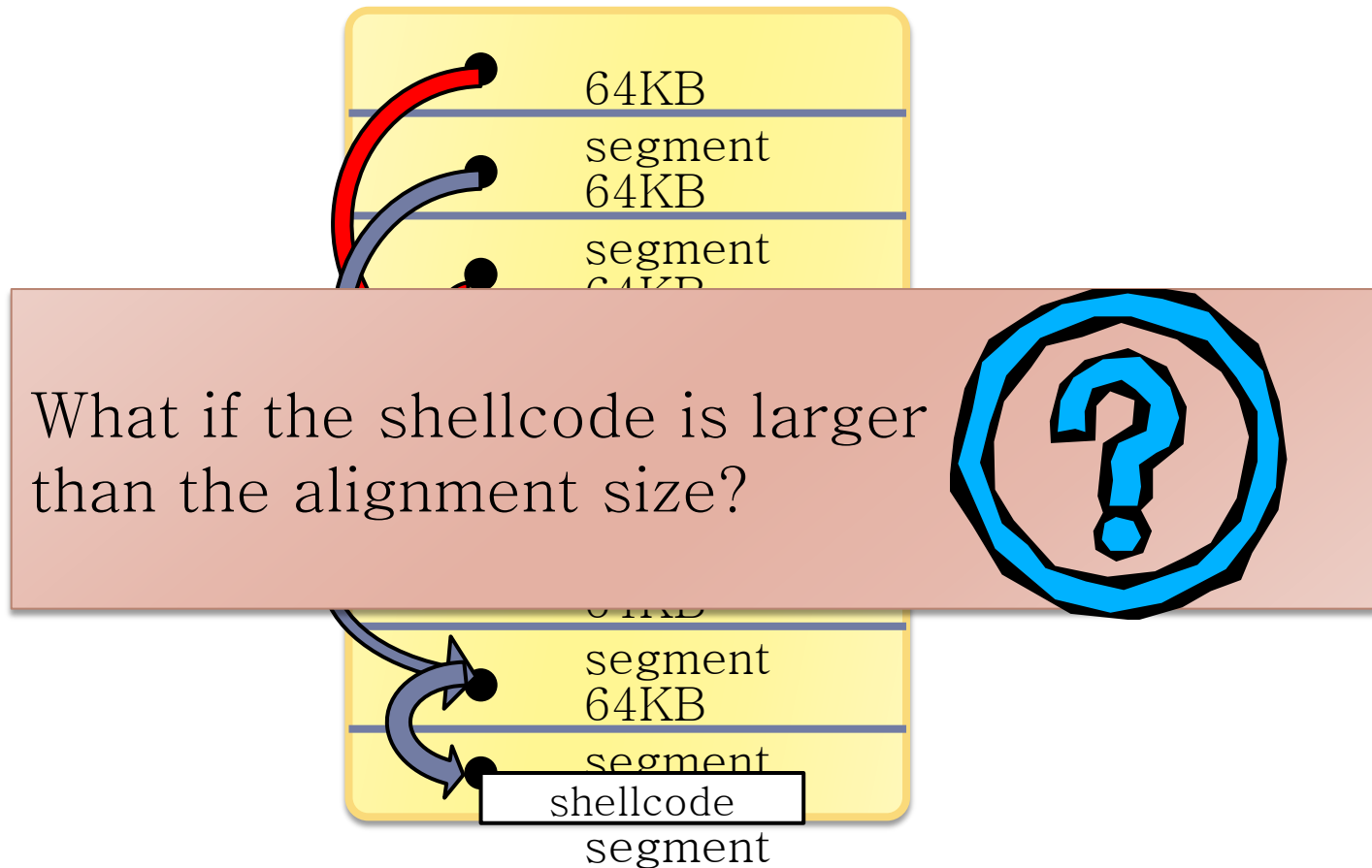
blocksize = 512Kbytes



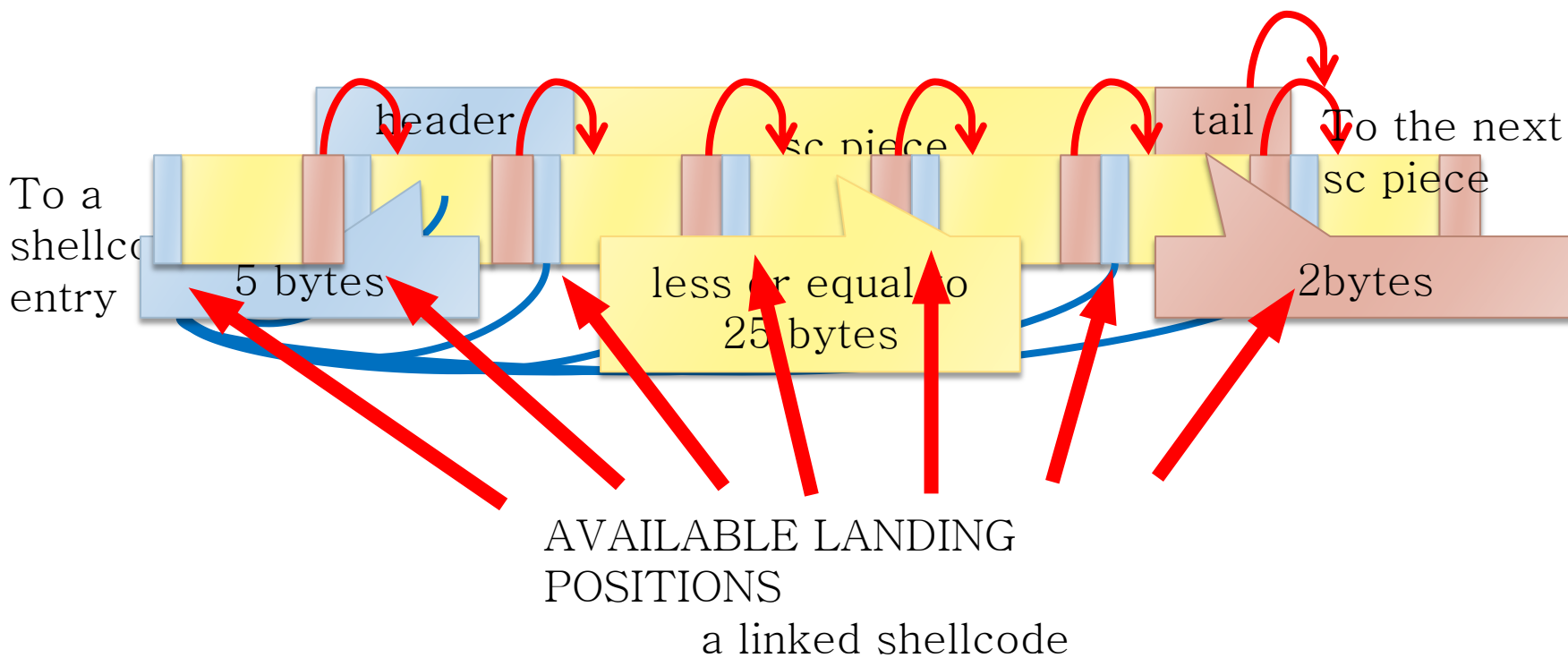
$$\text{NSA} \approx \text{sizeof}(\text{shellcode}) / 512\text{K} \approx 0\%$$



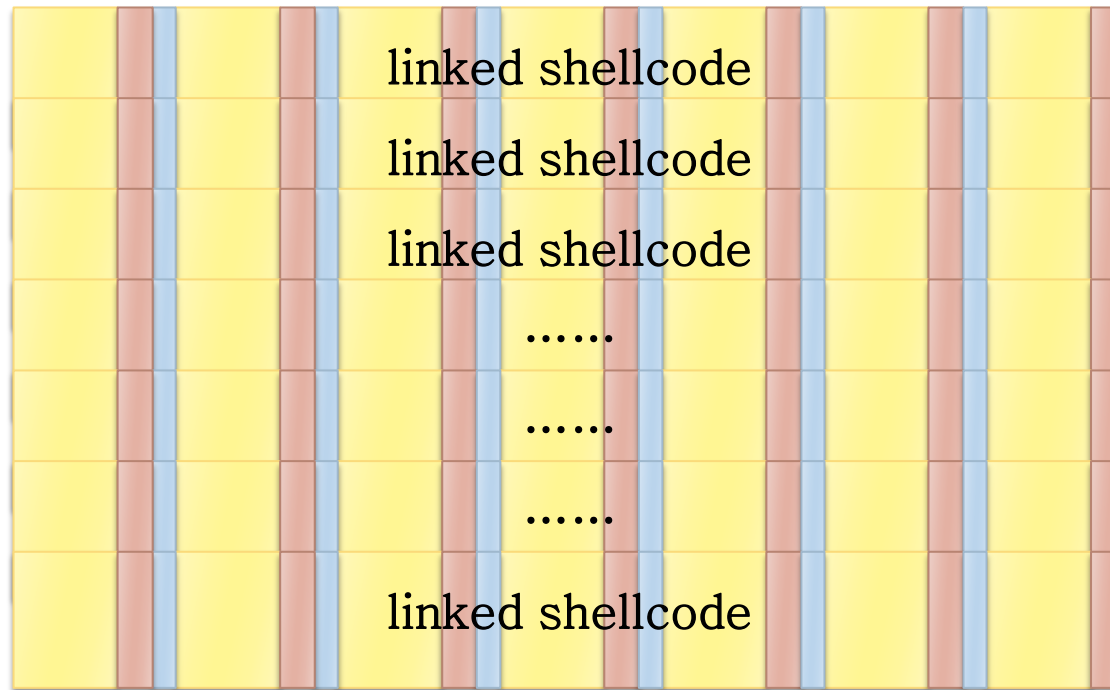
A Mix Type of TypeA, B, C



The TypeD Attack (alignment size = 32bytes)



The TypeD Attack (alignment size = 32bytes)



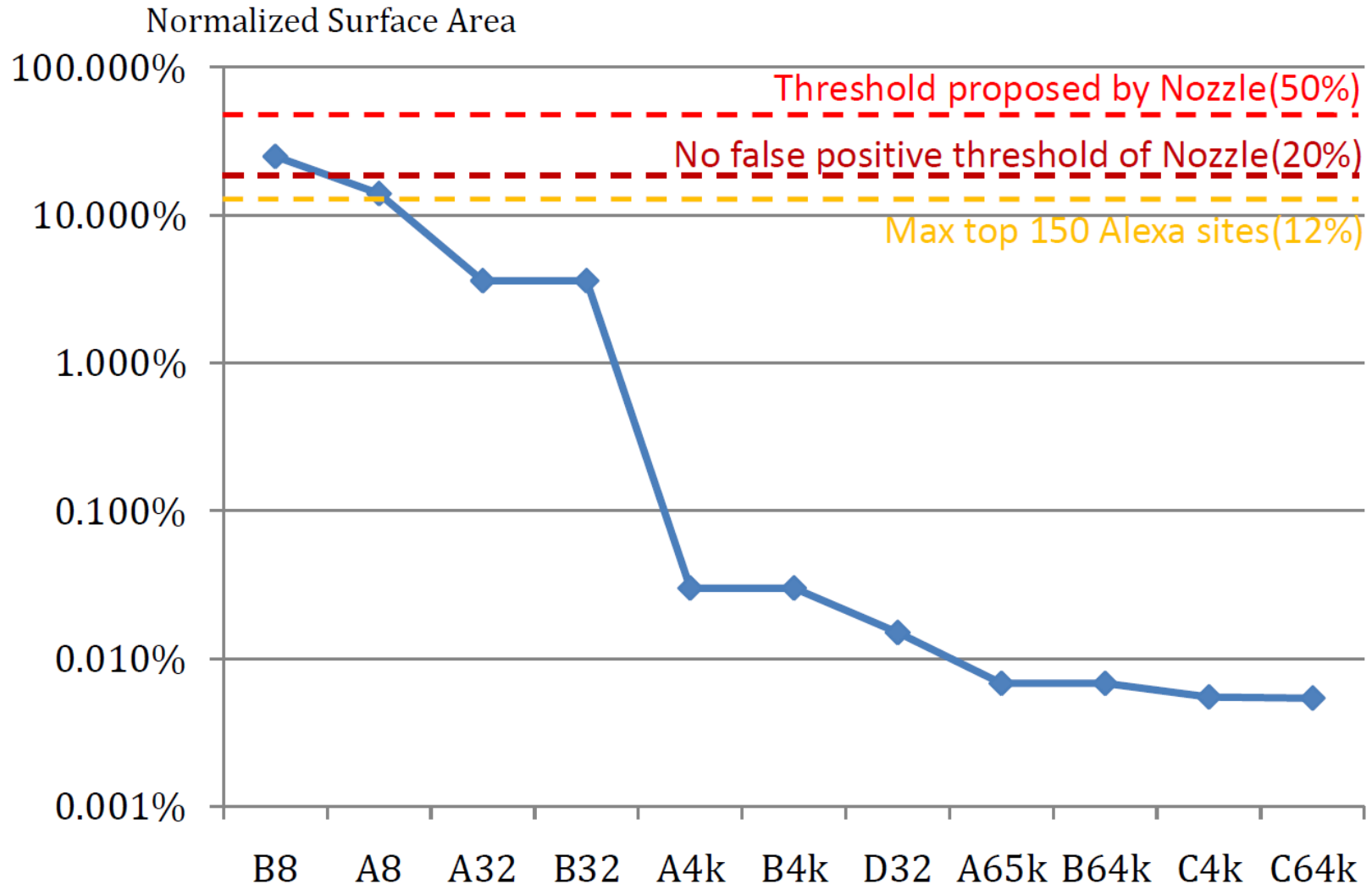
Relationship between Attack Type and Alignment Size

▶ Summarize

Alignment size	Type A	Type B	Type C	Type D
64 kbytes	✓	✓	✓	✓
32 bytes	✓	✓	X	✓
8 bytes	✓	✓	X	X
4 bytes	✓	X	X	X

- ▶ **Small amount of instructions keep high success rate with help of memory allocation granularity.**

Evaluation of Heap Taichi Attack Surface Area Calculation



Detecting Heap Taichi Attacks

Detecting Heap Taichi

Alignment size	Type A	Type B	Type C	Type D
64 kbytes	✓	✓	✓	✓
32 bytes	✓	✓	X	✓
8 bytes	✓	✓	✓	✓

1. Decrease the alignment size.
2. Improve the normalized surface area calculation.

Coarse granularity helps hide the shellcode and jump instructions.

using multiple copies of shellcode

Difficulty in Lower the Memory Allocation Granularity

Our Approach:
We modified JSarena and Firefox to support 8-byte alignment size

Application

JSarena in Firefox

alignment size
= 2Mbytes

Application

jemalloc

tcmalloc

.....

alignment size
= 8 bytes?

Kernel

heap manager
(Windows)

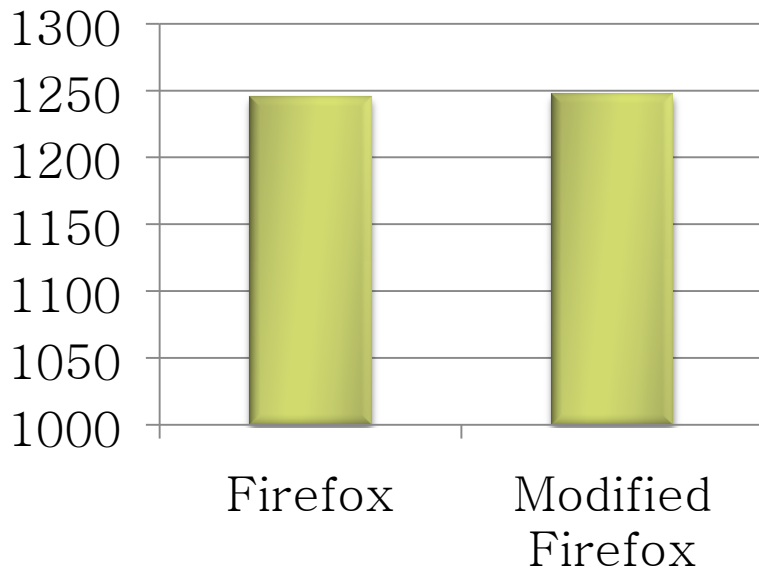
SLUB
allocator
(Linux)

.....

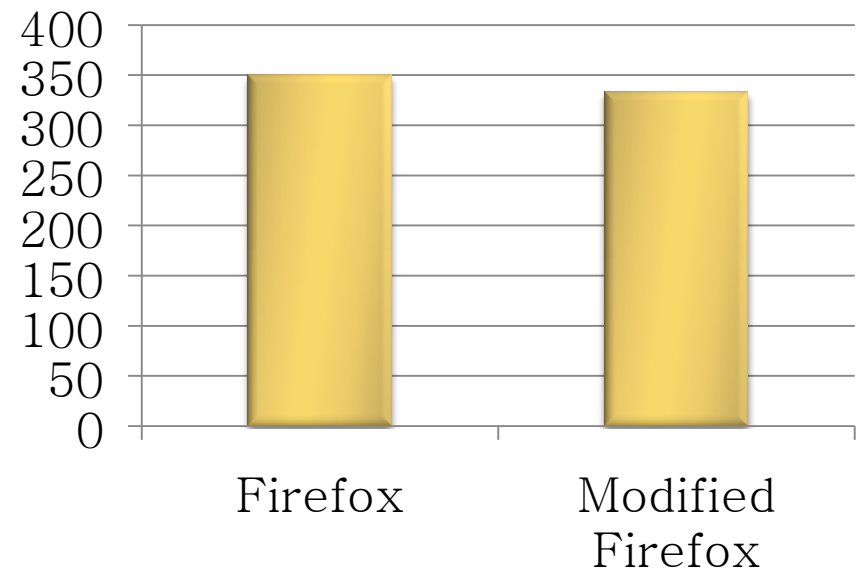
8-byte alignment is feasible

- ▶ The performance overhead is less than 5%

Sunspider Benchmark

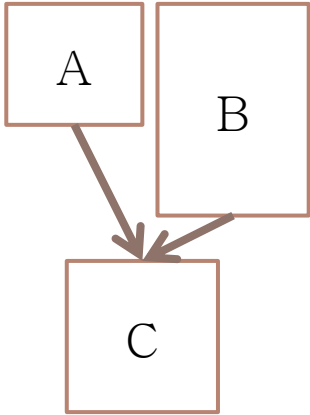


V8 Benchmark



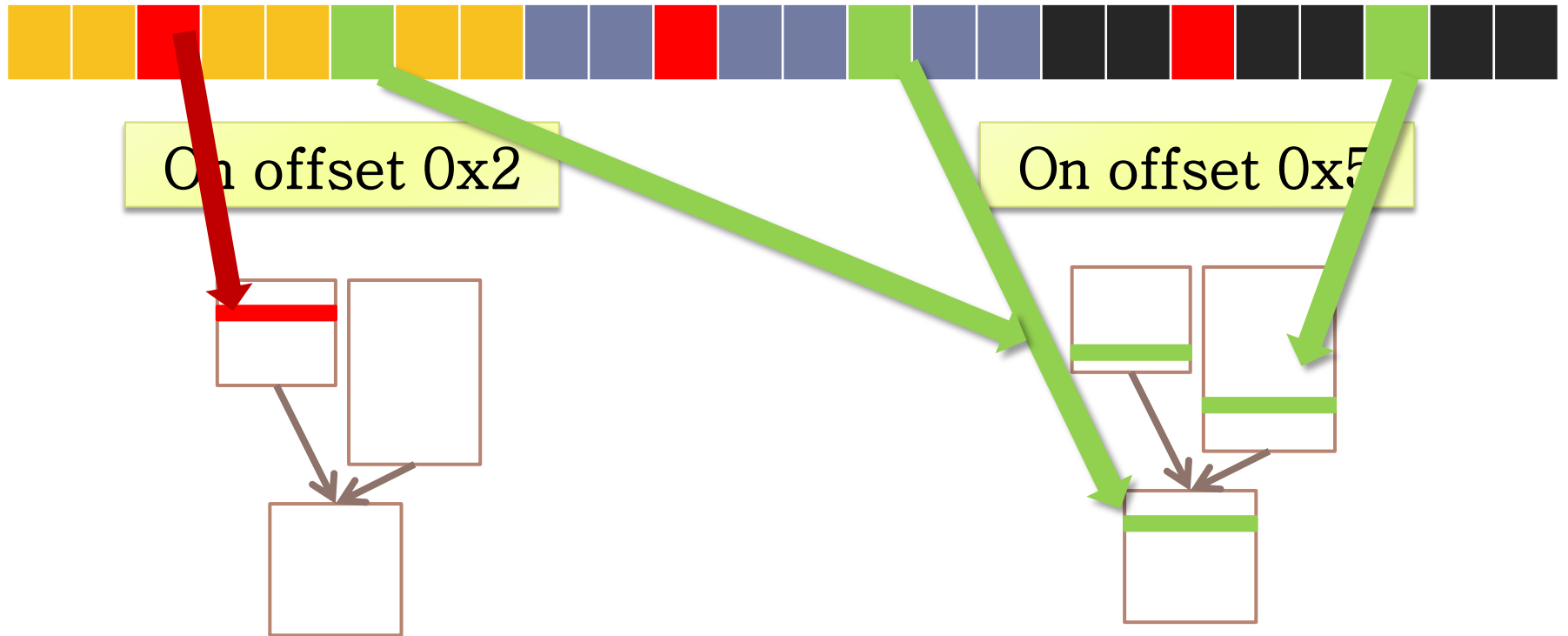
Enhance Nozzle

(Alignment size = 8 bytes)



Enhance Nozzle

(Alignment size = 8 bytes)



Success rate on offset 0x2
= $1/3 = 33\%$

RELATED WORK

- ▶ Shellcode-based detection
 - ▶ Egele et, al. DIMVA'09
 - ▶ Sam Small et, al. CCS'09
 - ▶ Yingbo Song et,al. CCS'07
- ▶ Heap spraying with ASLR
 - ▶ The granularity is still 64KBytes
- ▶ Heap spraying with DEP
 - ▶ DEP can be bypassed
 - ▶ Peter Vreugdenhil, Pwn2Own
- ▶ Solid Heap Memory Allocator
 - ▶ DieHard, DieHarder
- ▶ Memory exploit detection and prevention

Conclusion

- ▶ We analysis the weakness of memory alignment size.
- ▶ We present Heap Taichi, a new heap-spraying skill utilizing the weakness of memory alignments.
- ▶ We present an enhanced nozzle algorithm. The new algorithm can recognize these Heap Taichi attacks cooperating with finer memory allocation granularity.



Thank you!

